CATALOGUES OF PL-MANIFOLDS AND COMPLEXITY ESTIMATIONS VIA CRYSTALLIZATION THEORY

Maria Rita Casali Università di Modena e Reggio Emilia (Italy) casali@unimore.it

Workshop "TRIANGULATIONS" - Oberwolfach - May 4, 2012

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

Coloured triangulations and coloured graphs

A coloured triangulation of a compact PL *n*-manifold M^n is a pair (\bar{K}, ξ) , where \bar{K} is a pseudocomplex¹ triangulating M^n and $\xi : S_0(\bar{K}) \to \Delta_n = \{0, 1, \dots, n\}$ (vertex-labelling) satisfies:

- i) each *n*-simplex of \bar{K} has exactly one *c*-labelled vertex, for every $c \in \Delta_n$;
- ii) each *n*-labelled vertex is internal in \bar{K} .

EXAMPLE: If $M^n = |K|$, then (K', ξ) is a coloured triangulation, where

K' first barycentric subdivision of K $\xi(v) = r$ iff v barycenter of $\tau^r \in K$

A coloured triangulation \overline{K} of M^n is combinatorially visualized by means of an (n+1)-coloured graph (Γ, γ) :

- $\Gamma = (V(\Gamma), E(\Gamma))$ is the 1-skeleton of the dual cellular complex of \bar{K} ;
- $\gamma: E(\Gamma) \to \Delta_n$ (edge-coloration) is defined by: $\gamma(e) = c$ if $e \in E(\Gamma)$ is dual to an (n-1)-simplex of \overline{K} having no *c*-labelled vertex.

Г







Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

 (Γ, γ) is said to *represent* M^n , since the reversed process allows to completely reconstruct the coloured triangulation $\overline{K} = K(\Gamma)$ - and hence $M^n = |K(\Gamma)|$ - from it:

- 1) take an *n*-simplex $\sigma(x)$ for every vertex $x \in V(\Gamma)$, and label its vertices by Δ_n ;
- 2) if $x, y \in V(\Gamma)$ are joined by a *c*-coloured edge, identify the (n-1)-faces of $\sigma(x)$ and $\sigma(y)$ opposite to *c*-labelled vertices, so that equally labelled vertices coincide.

 (Γ, γ) is also said to be a *gem* ("graph <u>encoded manifold</u>") of M^n .



<ロ> <同> <同> < 回> < 回>

5 DQC

CONSEQUENCES:

- If Mⁿ is a closed manifold, any (n + 1)-coloured graph representing it is a regular graph of degree n + 1;
 If ∂Mⁿ ≠ Ø, any (n + 1)-coloured graph representing Mⁿ has a subset of vertices (boundary vertices) of degree n, lacking in n-coloured edges and corresponding to boundary n-simplices of K(Γ).
- $M^n = |K(\Gamma)|$ is orientable iff Γ is bipartite; " non-orientable " " " non-bipartite.
- ∀B ⊂ Δ_n, with #B = h, there is a bijection between (n − h)-simplices of K(Γ) whose vertices are labelled by Δ_n − {B} and connected components of h-coloured graph Γ_B = (V(Γ), γ⁻¹(B)).

A crystallization of an *n*-manifold M^n is any (n + 1)-coloured graph (Γ, γ) representing it, so that $K(\Gamma)$ has the <u>minimal</u> number of vertices.

If ∂M^n is either empty or connected, that minimal number is always equal to n + 1:

 (Γ, γ) is a crystallization of $M^n = |K(\Gamma)|$ if and only if $\Gamma_{\hat{c}}$ is connected, $\forall c \in \Delta_n$ (or, equivalently, if and only if $K(\Gamma)$ has exactly one c-labelled vertex, $\forall c \in \Delta_n$)

Pezzana Existence Theorem (1974)

Each PL *n*-manifold M^n (with or without boundary) admits a crystallization.



 \mathbb{S}^3

<ロ> <同> <同> < 回> < 回>

E 996



 $\mathbb{S}^1\times\mathbb{S}^2$

<ロ> <同> <同> < 回> < 回>

э.



 \mathbb{CP}^2

A complete (finite) set of graph-moves allows to translate the (PL)-homeomorphism problem for *n*-manifolds into an equivalence problem for (n + 1)-coloured graphs:

two coloured graphs represent the same PL-manifold if and only if they can be obtained one each other by a finite sequence of dipole moves.

An *h-dipole* $(1 \le h \le n)$ of (Γ, γ) is a subgraph $\Theta = \{v, w\}$ consisting of two vertices $v, w \in V(\Gamma)$ joined by *h* edges coloured by $j_1, j_2, \ldots, j_h \in \Delta_n$, such that:

- v and w belong to different components, Ξ_1 and Ξ_2 say, of $\Gamma_{\Delta_n \{j_1, \dots, j_h\}} = (V(\Gamma), \gamma^{-1}(\Delta_n \{j_1, \dots, j_h\}));$
- if either v or w is an internal vertex, then either Ξ_1 or Ξ_2 is a regular graph of degree n + 1 h.

Cataloguing PL-manifolds via crystallization theory Complexity estimations



dipole move

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

(日) (同) (三) (三)

2

2. Cataloguing PL-manifolds via crystallization theory

Each order 2p (n + 1)-coloured graph (Γ, γ) (with $V(\Gamma) = \{v_1, \dots, v_{2p}\}$) may be obviously encoded by an "incidence matrix"

$$A_{\Gamma}: \mathbb{N}_{2p} \times \mathbb{N}_{n+1} \rightarrow \{0, 1, \ldots, 2p\},\$$

$$A_{\Gamma}(i,c) = \begin{cases} j & \text{if } v_i \text{ is } c\text{-adjacent to } v_j \\ 0 & \text{if } v_i \text{ has no } c\text{-adjacent vertex} \end{cases}$$



The standard crystallization of $\mathbb{S}^2\times\mathbb{S}^1$

	colour 0	colour 1	colour 2	colour 3
vertex a	A	С	A	D
vertex A	а	b	а	d
vertex b	В	A	С	C
vertex B	b	с	d	с
vertex c	С	В	D	В
vertex C	с	а	b	b
vertex d	D	D	В	A
vertex D	d	d	С	а

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

æ

3

If we consider a manifold with boundary, zero elements appear, corresponding to boundary vertices of the graph.



A crystallization of $\mathbb{S}^1\times\mathbb{D}^2$

	colour 0	colour 1	colour 2	colour 3
vertex a	D	C	A	0
vertex A	d	b	а	0
vertex b	С	A	С	В
vertex B	с	с	d	b
vertex c	В	В	D	0
vertex C	b	а	b	0
vertex d	A	D	В	0
vertex D	а	d	с	0

э

The incidence matrix A_{Γ} is not the "most economical" way to identify (Γ, γ) (for example, $A(i, c) = j \iff A(j, c) = i$).

Moreover:

- if Γ is bipartite (i.e. it represents an orientable *n*-manifold), information about only one bipartition class is sufficient, for each colour c ∈ Δ_n;
- if Γ is non-bipartite (i.e. it represents a non-orientable *n*-manifold), for each colour c ∈ Δ_{n-1} information about only one bipartition class is sufficient, while adjacencies by colour n have to be completely described.

Finally, by suitably labelling the vertices of the $\Gamma,$ adjacencies by colour 0 may always be understood.



The standard crystallization of $\mathbb{S}^2\times\mathbb{S}^1$

	colour 0	colour 1	colour 2	colour 3
vertex a		С	A	D
vertex A				
vertex b		A	С	C
vertex B				
vertex c		В	D	В
vertex C				
vertex d		D	В	A
vertex D				

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

æ

The notion of CODE:

- for each vertex r ∈ V(Γ) and for every permutation ε of Δ_n = {0,1,...,n}, algorithmically and canonically label V(Γ) (so that the associated incidence matrix Ā^(r,ε)_Γ contains "essential" elements in well-defined positions);
- if $c_{r,\varepsilon}$ is the numerical string containing in orderly way the essential elements of $\bar{A}_{\Gamma}^{(r,\varepsilon)}$, the *code code*(Γ) is the lexicographic maximum among all strings $c_{r,\varepsilon}$:

$$code(\Gamma) = max \left\{ \begin{array}{l} r \in V(\Gamma) \\ \varepsilon \ permutation \ of \ \Delta_n \end{array} \right\}$$



The standard crystallization of $\mathbb{S}^2\times\mathbb{S}^1$

	$\varepsilon_0 = 0$	$\varepsilon_1 = 1$	$\varepsilon_2 = 3$	$\varepsilon_3 = 2$
vertex a		С	D	A
vertex A				
vertex b		А	С	C
vertex B				
vertex c		В	В	D
vertex C				
vertex d		D	А	В
vertex D				

$code(\Gamma) = CABD \ DCBA \ ACDB$

э

□ > 《注》《注》



The standard crystallization of $\mathbb{S}^2 \, \widetilde{\times} \, \mathbb{S}^1$

	$\varepsilon_0 = 0$	$arepsilon_1=1$	$\varepsilon_2 = 3$	$\varepsilon_3 = 2$
vertex a		С	D	A
vertex A				а
vertex b		A	С	С
vertex B				D'
vertex c		В	В	d′
vertex C				b
vertex d'=D		D	A	с
vertex D'=d				В

$$code(\Gamma) = CABD \ DCBA \ ACdc \ aDbB$$

æ

(1日) (1日) (1日)

The notion of code allows to detect *colour-isomorphic* graphs, i.e. graphs isomorphic up to permutation of the vertex set AND up to permutation of the colour set:

Theorem [C.-Gagliardi 2001]

 (Γ, γ) and (Γ', γ') are colour-isomorphic if and only if

 $code(\Gamma) = code(\Gamma').$

The notion of code is very useful in order to produce automatic catalogues of PL *n*-manifolds via crystallizations.

In the CLOSED case, it is necessary:

- to proceed inductively on dimension *n*;
- to perform sphere-recognition at every step

 (an (n + 1)-coloured graph (Γ, γ) represents an n-manifold if and
 only if Γ_c represents the (n − 1)-sphere S^{n−1}, ∀c ∈ Δ_n).

Advantages in dimension n = 3:

- (Γ, γ) is a crystallization of a 3-manifold M³ iff:
 i) Γ_ĉ is connected, ∀c ∈ Δ₃;
 - ii) $g_{01} + g_{02} + g_{03} = 2 + p$;
 - iii) $\forall \epsilon = (\epsilon_0, \epsilon_1, \epsilon_2, \epsilon_3), \ g_{\epsilon_0 \epsilon_1} = g_{\epsilon_2 \epsilon_3}.$
- all closed connected 3-manifolds may be represented by *rigid crystallizations*.

The generating algorithm was implemented in C++ programs starting from $\mathcal{S}^{(2p)}$ with $1 \leq p \leq 15$; the output data are presented in the following table according to the number of vertices.

2р	2	4	6	8	10	12	14	16	18	20	22	24	26	28	30
$\# \mathcal{C}^{(2p)}$	1	0	0	1	0	1	1	3	4	23	44	262	1252	7760	56912
$\# \tilde{C}^{(2p)}$	0	0	0	0	0	0	1	1	1	9	12	88	480	2790	21804

Table: rigid crystallizations up to 30 vertices

CLASSIFICATION ALGORITHM

After the generation process, suitable moves on gems, translating the PL-homeomorphism of the represented manifolds, are applied to develop a classification procedure which allows to detect crystallizations of the same manifold:

- dipole moves;
- generalized dipole moves (defined only for n = 3);
- switching of ρ-pairs (preserving the homeomorphism class, up to connected sum with handles).

The "admissible moves" are used to subdivide a given list of rigid crystallizations into equivalence classes, so that:

$$cl(\Gamma) = cl(\Gamma') \implies |K(\Gamma)| = |K(\Gamma')|$$

Remark: Note that we have no theoretical proof that

$$|K(\Gamma)| = |K(\Gamma')| \implies cl(\Gamma) = cl(\Gamma').$$

Nevertheless, by experimental results, the above implication is true for all elements of the catalogues $C^{(2p)}$, $\tilde{C}^{(2p)}$ (with $1 \le p \le 15$), with respect to a suitably chosen set \bar{S} of admissible moves.

In dimension 3 the above automatic partition into equivalence classes succeeds to distinguish topologically all manifolds represented by the generated catalogues:²

Proposition [C.-Cristofori 2008]

There exists a one-to-one correspondence between the set of classes of $\boldsymbol{C}^{(30)}$ (resp. $\tilde{\boldsymbol{C}}^{(30)}$) produced by the classification program and the set of orientable (resp. non-orientable) 3-manifolds admitting a coloured triangulation with at most 30 tetrahedra.

²For each positive integer p, we denote by $C^{(2p)}$ (resp. $\tilde{C}^{(2p)}$) the catalogue of all rigid bipartite (resp. non bipartite) crystallizations of order $\leq 2p$ arising from the generating algorithm.

Proposition [C.-Cristofori 2008]

There are exactly 110 closed prime orientable 3-manifolds, having a coloured triangulation with at most 30 tetrahedra. They are:

- fifty-five elliptic 3-manifolds;
- thirty-nine non-elliptic Seifert 3-manifolds (in particular, two torus bundles with Nil geometry);
- four torus bundles with Sol geometry;
- two manifolds of type (K² × I) ∪ (K² × I)/A (A ∈ GL(2; Z), det(A) = -1), with Sol geometry;
- seven non-geometric graph manifolds;
- three hyperbolic Dehn-fillings (of the complement of link 6_1^3).

The details both about the implementation of the classification algorithm (performed by the C++ program Γ -class) and about the obtained results, are available at the WEB page

http://cdm.unimo.it/home/matematica/casali.mariarita/CATALOGUES.htm

i.

The distribution of prime manifolds in $C^{(30)}$ with respect to *Matveev complexity* and *geometry*:

complexity	1	2	3	4	5	6	7	8	9	10
lens	2/2	3/3	6/6	10/10	0/20	0/36	0/72	0/136	0/272	0/528
other elliptic	-	1/1	1/1	4/4	11/11	14/25	0/45	0/78	0/142	0/270
\mathbb{E}^3	-	-	-	-	-	6/6	-	-	-	-
Nil	-	-	-	-	-	7/7	3/10	2/14	0/15	0/15
$\mathbb{H}^2\times\mathbb{R}$	-	-	-	-	-	-	-	0/2	-	0/8
$\widetilde{SL}_2(\mathbb{R})$	-	-	-	-	-	-	13/39	5/162	2/513	0/1416
Sol	-	-	-	-	-	-	4/5	2/9	0/23	0/39
non-geometric	-	-	-	-	-	-	4/4	1/35	2/185	0/777
hyperbolic	-	-	-	-	-	-	-	-	2/4	1/25
TOTAL	2/2	4/4	7/7	14/14	11/31	27/74	24/175	10/436	6/1154	1/3078

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

3 x 3

Proposition [C. 1999] - [Bandieri-Cristofori-Gagliardi 2009]

There are exactly sixteen closed prime non-orientable 3-manifolds, having a coloured triangulation with at most 30 tetrahedra. They are:

- the 3-manifolds $\mathbb{S}^1 \widetilde{\times} \mathbb{S}^2$ and $\mathbb{RP}^2 \times \mathbb{S}^1$, with $\mathbb{S}^2 \times \mathbb{R}$ geometry;
- the four non-orientable euclidean 3-manifolds;
- the torus bundles $TB\begin{pmatrix} 0 & 1 \\ 1 & -1 \end{pmatrix}$, $TB\begin{pmatrix} 2 & 1 \\ 1 & 0 \end{pmatrix}$, $TB\begin{pmatrix} 3 & 2 \\ 2 & 1 \end{pmatrix}$ and $TB\begin{pmatrix} 3 & 1 \\ 1 & 0 \end{pmatrix}$, with Sol geometry;
- the Seifert fibered spaces $(\mathbb{RP}^2; (2, 1), (3, 1)), (\overline{D}; (2, 1), (3, 1)), (T^2/o_2; (2, 1)), (K^2; (2, 1)) and (K^2/n_3; (2, 1)), with <math>\mathbb{H}^2 \times \mathbb{R}$ geometry;
- the non-geometric graph manifold

$$(\mathbb{A};(2,1))\cup(\mathbb{A};(2,1))/\begin{pmatrix}0&-1\\1&0\end{pmatrix}.$$

The distribution of prime manifolds in $\tilde{\textbf{C}}^{(30)}$ with respect to Matveev complexity and geometry:

complexity	0	1	2	3	4	5	6	7	8	9
$\mathbb{S}^2 imes \mathbb{R}$	1/1	1/1	-	-	-	-	-	-	-	-
\mathbb{E}^3	-	-	-	-	-	-	4/4	-	-	-
$\mathbb{H}^2 imes\mathbb{R}$	-	-	-	-	-	-	-	2/2	0/8	3/25
Sol	-	-	-	-	-	-	1/1	1/1	2/2	0/2
non-geometric	-	-	-	-	-	-	-	-	-	1/6
TOTAL	1/1	1/1	0	0	0	0	5/5	3/3	2/10	4/33

э

Two ideas for improving the generation process:

- Simplify the considered catalogues, by means of additional conditions on the representing objects (for example by considering only **rigid cluster-less crystallizations**, which represent ALL 3-manifolds).
- Improve the implementation by making use of a refined **parallel** version of the generating algorithm.

Preliminary problem

In dimension 4, the generation of crystallization catalogues implies the previous generation of all gems (not necessarily crystallizations!) representing 3-dimensional spheres up to a fixed order.

For the first segment of the catalogue, the recognition of gems representing \mathbb{S}^3 is easily faced and solved by dipole eliminations, since:

no rigid crystallization of S³ exists (different from the "trivial" one, with order two) with less than 24 vertices.

The generation choice

In order to avoid the explosion of data, in dimension 4 a great attention must be paid to the choice of the representing set.

Proposition [Bandieri-Gagliardi 2011]

Each closed connected PL 4-manifold M admits a *rigid* crystallization. Moreover, if M is handle-free, it admits a rigid crystallization of minimal order.

As a consequence:

- the catalogue of crystallizations representing closed 4-manifolds may be restricted to **rigid crystallizations lacking in** 2-**dipoles**;
- the catalogue of (non-contracted) gems representing S³ and lacking in ρ₃-pairs may be considered as input data of the 4-dimensional generating process.

The first output of the generating program:

2р	2	4	6	8	10	12	14	16	18	20
Γ ₄ (S³)	1	0	0	9	39	400	5255	95870	1994952	45654630
$\Gamma(M^4)$ lacking in 2-dipoles	1	0	0	1	0	0	1109	4512	44803	47623129

Table: number of gems of \mathbb{S}^3 lacking in ρ_3 -pairs and number of rigid crystallizations of 4-manifolds lacking in 2-dipoles, according to vertex number

Remark: A 4-dimensional crystallization catalogue - where classification is performed up to PL-homeomorphism - might provide examples of different PL 4-manifolds belonging to the same TOP-homeomorphism class.

The *gem-complexity* of a closed *n*-manifold M^n is the non-negative integer $k(M^n) = p - 1$, 2p being the minimum order of a crystallization of M^n .

Crystallization theory easily implies:

 $k(M^4) \geq 3\beta_2(M^4)$ for each simply connected closed 4-manifold M^4 .

Relation $k(M^4) \ge 3\beta_2(M^4)$, combined with the up-to-date topological classification of simply connected PL 4-manifolds allows to prove:

Proposition (to appear)

If M^4 is a simply connected closed PL 4-manifold with $k(M^4) \le 65$, then M^4 is TOP-homeomorphic to either

$$(\#_r \mathbb{CP}^2) \# (\#_{r'} - \mathbb{CP}^2)$$
 with $r + r' = \beta_2(M^4)$

or

$$\#_s(\mathbb{S}^2 \times \mathbb{S}^2)$$
 with $s = \frac{1}{2}\beta_2(M^4)$

In fact, the important results by

- [Freedman, *The topology of four-dimensional manifolds*, J. Differential Geom. 17 (1982), 357-453]
- [Donaldson, An application of gauge theory to four-dimensional topology, J. Differential Geom. 18 (1983), 279-315]
- [Furuta, Monopole equation and the $\frac{11}{8}$ conjecture, Math. Res. Lett. 8 (2001), 279-291]

ensure that only intersection forms of type

$$r[1] \oplus r'[-1]$$
 or $s \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$

are allowed if $k(M^4) \leq 65$, since forms of type

$$\pm 2nE_8 \oplus s \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

are proved to represent a PL 4-manifold only if s > 2n (and hence, only PL 4-manifolds with $\beta_2 \ge 22$ occur in this case).

RESULTS BY THE FIRST SEGMENT OF 4-DIMENSIONAL CATALOGUES:

Proposition (to appear)

- S⁴ is the only closed connected (PL) 4-manifold with gem-complexity 0.
- No closed connected (PL) 4-manifold M^4 exists with $1 \le k(M^4) \le 2$.
- The complex projective plane \mathbb{CP}^2 is the only closed connected (PL) 4-manifold with $k(M^4) = 3$.
- k(S¹ × S³) = k(S¹×S³) = 4; moreover, no closed connected handle-free (PL) 4-manifold M⁴ exists with k(M⁴) = 4.
- No closed connected (PL) 4-manifold M^4 exists with $k(M^4) = 5$.

RESULTS BY 4-DIMENSIONAL CATALOGUES (UP TO 18 VERTICES): THE NON-ORIENTABLE CASE

Proposition (to appear)

- No closed connected handle-free non-orientable (PL) 4-manifold M⁴ exists with k(M⁴) ≤ 6.
- The real projective space ℝP⁴ is the only closed connected prime non-orientable (PL) 4-manifold M⁴ with k(M⁴) = 7.
- No closed connected handle-free non-orientable (PL) 4-manifold M⁴ exists with 8 ≤ k(M⁴) ≤ 9.

RESULTS BY THE 4-DIMENSIONAL CATALOGUES (FROM 14 TO 18 VERTICES): THE ORIENTABLE CASE

Proposition (to appear)

- $k(\mathbb{S}^2 \times \mathbb{S}^2) = 6$; moreover, $k(\mathbb{CP}^2 \# \mathbb{CP}^2) = k(\mathbb{CP}^2 \# (-\mathbb{CP}^2)) = 6$, too.
- If M^4 is a closed connected handle-free orientable (PL) 4-manifold with $6 \le k(M^4) \le 9$,) then M^4 is simply-connected and TOP-homeomorphic to either $\mathbb{S}^2 \times \mathbb{S}^2$ or $\mathbb{CP}^2 \#\mathbb{CP}^2$ or $\mathbb{CP}^2 \#(-\mathbb{CP}^2)$.

3. Complexity estimations

By making use of the strong connection existing in dimension 3 between gems and Heegaard diagrams, a 3-manifold invariant based on crystallization theory - called *GM-complexity* - has been introduced and proved to be an upper bound for the Matveev complexity of each compact 3-manifold.

- M.R. C., Computing Matveev's complexity of non-orientable 3-manifolds via crystallization theory, Topology Appl. 144 (2004), 201-209.
- M.R. C., Estimating Matveev's complexity via crystallization theory, Discrete Math. 307 (2007), 704-714.
- M.R. C. P. Cristofori, Computing Matveev's complexity via crystallization theory: the orientable case, Acta Appl. Math. 92 (2006), 113-123.
- M.R. C. P. Cristofori M. Mulazzani, Complexity computation for compact 3-manifolds via crystallizations and Heegaard diagrams, Topology and its Applications (2012), to appear.
- M.R. C. P. Cristofori, Computing Matveev's complexity via crystallization theory: the boundary case, preprint 2012.

4 A N 4 B N 4 B N

3. Estimating Matveev complexity via Heegaard diagrams

If $\mathcal{H} = (S, v, w)$ is a Heegaard diagram of M, then a special spine of M exists, whose true vertices are the intersection points of the curves of the two systems v and w, with the exception of those lying on the boundary of a region of $S - \{v \cup w\}$.

Hence:

 $c(M) \leq n - m$,

where n = number of intersection points between v and w and m = number of intersection points contained in \overline{R} .



- H_{α} = the largest 2-dimensional subcomplex of $K'(\Gamma)$ not intersecting the barycentric subdivisions of $K_{\alpha 3}$ e $K_{\beta\beta'}$.

•
$$F_{\alpha} = |H_{\alpha}|.$$

The surface F_{α} splits $K(\Gamma)$ into two polyhedra $\mathcal{A}_{\alpha,3}$ and $\mathcal{A}_{\alpha',\beta'}$, whose intersection is exactly F_{α} :

- $\mathcal{A}_{\alpha 3} \searrow \mathcal{K}_{\alpha 3}$ (resp. $\mathcal{A}_{\beta \beta'} \searrow \mathcal{K}_{\beta \beta'}$) \Longrightarrow is an handlebody.
- $\mathcal{A}_{\alpha 3} \cap \mathcal{A}_{\beta \beta'} = \partial \mathcal{A}_{\alpha 3} \cap \partial \mathcal{A}_{\beta \beta'} = F_{\alpha}.$
- If Γ is a crystallization, let D be a $\{\beta, \beta'\}$ -coloured cycle and D' a $\{\alpha, 3\}$ -coloured cycle.

 $(F_{\alpha}, \Gamma_{\beta\beta'} \setminus D, \Gamma_{\alpha3} \setminus D')$ is a Heegaard diagram of M.

Example: Poincarè homology sphere



Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

э

GM-complexity

GM-complexity of a crystallization

$$c_{GM}(\Gamma) = min\{\#V(\Gamma) - \#(V(D) \cup V(D') \cup V(\Xi)) / \alpha \in \Delta_3, \\ D \in \Gamma_{\beta\beta'}, \ D' \in \Gamma_{\alpha3}, \ \Xi \in \mathcal{R}_{\alpha}(D, D')\}$$

GM-complexity of a closed 3-manifold

 $c_{GM}(M) = min\{c_{GM}(\Gamma) \mid \Gamma \text{ crystallization of } M\}$

э

∃ ▶ ∢

Remark:

The computation of $c_{GM}(\Gamma)$ is only based on the combinatorial structure of the edge-coloured graph Γ ; hence, it may be determined in an algorithmic way (*GM-COMPLEXITY program*).

For details, see:

http://cdm.unimo.it/home/matematica/casali.mariarita/about_cgm.htm

By making use of this program, an estimation of c(M) has been obtained for all manifolds involved in existing crystallization catalogues.

Conjecture

For every closed connected 3-manifold M,

$$c(M) = c_{GM}(M).$$

GM-complexity: non-contracted case

GM-complexity of a gem

$$c_{GM}(\Gamma) = \min\{\#V(\Gamma) - \#(V(\mathcal{D}) \cup V(\mathcal{D}') \cup V(\Xi)) / \\ \alpha \in \Delta_2, \\ \mathcal{D} \text{ set of } \{\beta, \beta'\} - \text{ cycles dual to a maximal tree of } K_{\alpha 3}, \\ \mathcal{D}' \text{ set of } \{\alpha, 3\} - \text{ cycles dual to a maximal tree of } K_{\beta \beta'}, \\ \Xi \in \mathcal{R}_{\alpha}(\mathcal{D}, \mathcal{D}')\}$$

э

Complexity estimations [C. 2007]

The notion of *GM*-complexity, combined with the widely investigated relationships between crystallization theory and other representation methods for 3-manifolds, has allowed to obtain direct estimations of the Matveev complexity for several classes of manifolds, significantly improving former results.

Proposition. Let $M_2(L)$ be the (unique) 2-fold covering of \mathbb{S}^3 branched on link L and let \overline{L} be a p-bridge projection of L with k crossing points $(k \ge p \ge 2)$. If \overline{L} admits a bridge $\overline{\beta}$ with order m_b and an independent arc $\overline{\alpha}$ with relative order m_a (i.e. an arc not consecutive to $\overline{\beta}$ and containing $m_a \ge 0$ crossing points not belonging to $\overline{\beta}$), then

$$c(M_2(L)) \leq 2(k + p - m_b - m_a - 3) \leq 4k - 8.$$



[M. Ferri, *Crystallisations of 2-fold branched coverings of* \mathbb{S}^3 , Proc. Amer. Math. Soc. **73** (1979), 271-276.]

э

∃ → < ∃ →</p>

____ ▶

Complexity estimations [C. 2007]

Proposition. Let $M^3(K, \omega)$ be the 3-fold simple covering of \mathbb{S}^3 branched on knot K with monodromy ω and let (\bar{K}, ω) be a 3-coloured diagram of knot K associated to the pair (K, ω) , with k crossing points and an order m arc. Then

$$c(M^{3}(K,\omega)) \leq 2k - 2(m+2).$$



[M.R. C., Coloured knots and coloured graphs representing 3-fold simple coverings of S^3 , Discrete Math. **137** (1995), 87-98.]

э

- 4 回 ト - 4 回 ト

Complexity estimations [C. 2007]

Proposition. Let $M^3(L, d)$ be the 3-manifold obtained by Dehn surgery on the framed link (L, d) and let \overline{L} be a planar (connected) diagram of Lwith $l \ge 1$ components and k crossing points. If \overline{L} admits a region of order m whose boundary involves components L_{j_1}, \ldots, L_{j_s} $(1 \le s \le l)$ of L, then

$$c(M^{3}(L,d)) \leq 6k + 2t - 4l - 2(m-1) - \sum_{p=1,...,s} t_{j_{p}},$$

where $t = \sum_{i=1,\ldots,l} t_i$ and $t_i = |d_i - w(\overline{L}_i)|, \forall i = 1,\ldots,l$.



[M.R. C., From framed links to crystallizations of bounded 4-manifolds, J. Knot Theory Ramifications **9** (2000), 443-458.]

A B > A B >

____ ▶

3

case $\partial M \neq \emptyset$



- $\partial \mathcal{A}_{\alpha 3} \cap \partial \mathcal{A}_{\beta \beta'} = F_{\alpha}$
- $\partial \mathcal{A}_{\alpha 3} \cap \partial M = \bigcup_{i=1}^{r} \mathbb{D}_{i}$
- $\partial \mathbb{D}_i = lk(v_i, (\partial K)')$ $v_i \quad \alpha - labelled vertex of <math>\partial K$ $\bigcup_{i=1}^r \partial \mathbb{D}_i = \partial F_{\alpha}$

•
$$S_{\alpha} = F_{\alpha} \cup (\bigcup_{i=1}^{r} \mathbb{D}_{i})$$

•
$$C = S_{\alpha} \times [-1, +1]$$
 collar of
 S_{α} in $\mathcal{A}_{\alpha 3}$
 $C^{-} = S_{\alpha} \times [-1, 0]$
 $C^{+} = S_{\alpha} \times [0, +1]$
 $C_{i} = S_{\alpha} \times \{i\}$

$$X_{\alpha} = \overline{\mathcal{A}_{\alpha 3} \setminus C^{-}}$$
$$Y_{\alpha} = \mathcal{A}_{\beta \beta'} \cup C^{-}$$

э

- X_α is an handlebody obtained from C⁺ by attaching 2-handles to C₁ along the {β, β'}-coloured cycles of Γ (dual to the edges of K_{α3}).
- Y_{α} is a compression body³, obtained from C^- by attaching 2-handles to C_{-1} along the $\{\alpha, 3\}$ -coloured cycles of Γ (dual to edges of $\mathcal{K}_{\beta\beta'}$ not belonging to $\partial \mathcal{K}$).

 $(S_{\alpha}, X_{\alpha}, Y_{\alpha})$ is a (generalized) Heegaard splitting of M.

³A compression body is a 3-manifold with boundary obtained from a $F \times I$ (F surface) by attaching 2-handles and 3-handles to $F \times \{1\}$. $\Box \mapsto \langle \Box \rangle \land \langle \Xi \rangle \land \langle \Xi \rangle$

GM-complexity of a gem with boundary

- (a) $\mathcal{D} = \text{set of } \{\beta, \beta'\}$ coloured cycles dual to a maximal tree of $K_{\alpha 3}$.
- (b) $G_{\beta\beta'}$ = graph obtained from $K_{\beta\beta'}$ by contracting to one point p_i (for each i = 1, ..., r) the vertices of $K_{\beta\beta'}$ belonging to the *i*-th component of ∂K .
- (c) $\mathcal{D}' = \text{set of } \{\alpha, 3\}$ -coloured cycles dual to the edges of a subgraph \overline{G} of $G_{\beta\beta'}$ such that \overline{G} is a union of trees containing all vertices of $G_{\beta\beta'}$ and, $\forall i, j, i \neq j$, the vertices p_i and p_j belong to different connected components of \overline{G} .

GM-complexity of a gem with boundary

$$c_{GM}(\Gamma) = \min\{\#V(\Gamma) - \#(V(\mathcal{D}) \cup V(\mathcal{D}') \cup V(\Xi)) / \\ \alpha \in \Delta_2, \\ \mathcal{D}, \mathcal{D}' \text{ satisfying (a) and (c)} \\ \Xi \in \mathcal{R}_{\alpha}(\mathcal{D}, \mathcal{D}')\}$$

GM-complexity of a 3-manifold with boundary

GM-complexity of a 3-manifold with boundary

 $c_{GM}(M) = \min\{c_{GM}(\Gamma) \mid \Gamma \text{ gem of } M\}$

By definition itself:

Proposition [C.-Cristofori 2012]

For each compact 3-manifold M,

 $c(M) \leq c_{GM}(M)$

Maria Rita Casali UP-TO-DATE RESULTS IN CRYSTALLIZATION THEORY

WORK IN PROGRESS: estimations via *c*_{GM} (boundary case)

By making use of the notion of *GM-complexity* for 3-manifolds with boundary, we hope to obtain - via graph theoretical construction of the associated gems - improvements for existing estimations of Matveev complexity for some interesting classes of bounded 3-manifolds.

The first efforts will take into considerations

knot (or link) complements.